

**$m$ -Hypergeometric Solutions of Anti-Difference and  
 $q$ -Anti-Difference Equations**

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**Abstract**

In this paper we consider the problem of finding  $m$ -hypergeometric solutions of anti-difference equations. We extend the greatest factorial factorization (GFF) of a polynomial, introduced by Paule (1995), to the  $m$ -greatest factorial factorization ( $m$  GFF). Equipped with the  $m$  GFF-concept, we present algebraically motivated approach to the problem. This approach requires only “gcd” operations but no factorization. Then, we solve the same problem for  $q$ -anti-difference equations.

*Keywords* : Gosper’s algorithm,  $m$ -hypergeometric solution,  $m$ -greatest factorial factorization,  $q$ -Gosper algorithm,  $qm$ -hypergeometric solution,  $qm$ -greatest factorial factorization.

$q$ -

$m$ -  
m- (1995) (GFF)  
gcd (mGFF)  
 $q$ -

## 1. Introduction

Let  $m$  denotes a positive integer,  $\mathbb{N}$  be the set of natural numbers,  $K$  be the field of characteristic zero,  $K(n)$  be the field of rational functions over  $K$ ,  $K[n]$  be the ring of polynomials over  $K$ ,  $F$  denotes the transcendental extension of  $K$  by the indeterminate  $q$ , i.e.,  $F = K(q)$ ,  $E$  denotes the shift operator on  $K[n]$ , i.e.,  $(Ep)(n) = p(n+1)$  for any  $p \in K[n]$ ,  $\varepsilon$  denotes the  $q$ -shift operator on  $F[n]$  and  $F(n)$ , i.e.,  $(\varepsilon u)(n) = u(qn)$  for any  $u \in F[n]$  or  $u \in F(n)$ ,  $\deg(p)$  denotes the polynomial degree (in  $n$ ) of any  $p \in K[n]$  or  $p \in F[n]$ ,  $p \neq 0$ . We define  $\deg(0) = -1$ . We assume the result of any gcd (greatest common divisor) computation in  $K[n]$  or  $F[n]$  as being normalized to a monic polynomial  $p$ , i.e., the leading coefficient of  $p$  being 1. Recall that a non-zero term  $t_n$  is called a hypergeometric term over  $K$  if there exist a rational function  $r(n) \in K(n)$  such that

$$\frac{t_{n+1}}{t_n} = r(n).$$

Gosper's algorithm (Gosper, 1978) (also see Graham *et al.*, 1989, Koepf, 1998, Petkovšek *et al.*, 1996) has been extensively studied and widely used to prove hypergeometric identities. Given a hypergeometric term  $t_n$ , Gosper's algorithm is a procedure to find a hypergeometric term  $z_n$  satisfying

$$z_{n+1} - z_n = t_n. \quad (1.1)$$

if it exists, or confirm the nonexistence of any solution of (1.1). In Paule (1995), Paule introduced the GFF-concept. Equipped with the GFF-concept, he presented a new and algebraically motivated approach to Gosper's algorithm.

A non-zero term  $a_n$  is called an  $m$ -hypergeometric over  $K$  if there exist a rational function  $w(n) \in K[n]$  such that

$$\frac{a_{n+m}}{a_n} = w(n). \quad (1.2)$$

In Koepf (1995), Koepf extends Gosper's algorithm to find  $m$ -hypergeometric solutions  $h_n$  of

$$h_{n+m} - h_n = a_n, \quad (1.3)$$

where  $a_n$  is a given  $m$ -hypergeometric term. In Petkovšek and Bruno (1993), Petkovšek and Bruno described an algorithm to find  $m$ -hypergeometric solutions of homogeneous linear recurrences with polynomial coefficients. Their algorithm reduces to algorithm **Hyper** (Petkovšek, 1992) when  $m = 1$ .

A non-zero term  $b_k$  is called a  $q$ -hypergeometric over  $F$  if there exists a rational function  $\sigma \in F(q^k)$  such that

$$\frac{b_{k+1}}{b_k} = \sigma(q^k).$$

In Paule and Riese (1997), Paule and Riese introduced the  $q$ -greatest factorial factorization ( $q$  GFF) of polynomials, which is a  $q$ -analogue of the GFF-concept. Equipped with the  $q$  GFF, they presented a new approach to find  $q$ -hypergeometric solutions  $l_k$  of

$$l_{k+1} - l_k = b_k, \quad (1.4)$$

where  $b_k$  is a given  $q$ -hypergeometric term. Paule-Riese's approach can be viewed as an  $q$ -analogue of Gosper's algorithm.

A non-zero term  $f_k$  is called a  $qm$ -hypergeometric over  $F$  if there exist a rational function  $\rho \in F(q^k)$  such that

$$\frac{f_{k+m}}{f_k} = \rho(q^k).$$

Let us define the dispersion  $\text{dis}_m(a, b)$  of the polynomials  $a(n), b(n) \in K[n]$  to be the greatest nonnegative integer  $k$  (if it exists) such that  $a(n)$  and  $b(n + mk)$  have a nontrivial common divisor, i.e.,

$$\text{dis}_m(a, b) = \max \{ k \in \mathbb{N} \mid \deg \gcd(a(n), b(n + mk)) \geq 1 \}.$$

If  $k$  does not exist then we set  $\text{dis}_m(a, b) = -1$ . Recall that the pair  $\langle c, d \rangle$ ,  $c, d \in K[n]$ , is called the reduced form of  $r \in K(n)$  if  $r = \frac{c}{d}$ ,  $d$  is monic, and  $\gcd(c, d) = 1$ .

The contents of this paper are as follows: In Section 2, we give the Fundamental  $m$  GFF Lemma, which is an extension of the Fundamental Lemma given by Paule (1995). In Section 3, we extend Paule's approach to find  $m$ -hypergeometric solutions of anti-difference equations. In Section 4, we give the Fundamental  $qm$  GFF Lemma, which is an extension of the Fundamental  $q$  GFF Lemma given by Paule and Riese (1997). Finally, In Section 5, we extend Paule-Riese's approach to find  $qm$ -hypergeometric solutions of  $q$ -anti-difference equations.

## 2. $m$ -Greatest Factorial Factorization

In this section we define the  $m$  GFF of a polynomial, which is an extension of the GFF-concept introduced by Paule.

## 2.1 Basic Definitions

**Definition 2.1.** For any monic polynomial  $p \in K[n]$  and  $i \in \mathbb{N}$ , the  $i$ -th  $m$ -falling factorial  $[p]_m^i$  of  $p$  is defined as

$$[p]_m^i = p \cdot E^{-m} p \cdot E^{-2m} p \cdot \dots \cdot E^{-(i+1)m} p.$$

For  $i = 0$ , we let  $[p]_m^0 = 1$ .

**Definition 2.2.** We say that  $\langle p_1, p_2, \dots, p_s \rangle$ ,  $p_i \in K[n]$ , is an  $m$  GFF-form of a monic polynomial  $p \in K[n]$  if the following conditions hold:

- ( $m$  GFF1)  $p = [p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_s]_m^s$ ,
- ( $m$  GFF2) each  $p_i$  is monic and  $s > 0$  implies  $\deg(p_s) > 0$ ,
- ( $m$  GFF3)  $\gcd([p_i]_m^i, E^m p_j) = 1$  for  $1 \leq i \leq j \leq s$ ,
- ( $m$  GFF4)  $\gcd([p_i]_m^i, E^{-jm} p_j) = 1$  for  $1 \leq i \leq j \leq s$ .

If  $\langle p_1, p_2, \dots, p_s \rangle$  is an  $m$  GFF-form of a monic  $p \in K[n]$  we sometimes express this fact for short by  $m$  GFF( $p$ ) =  $\langle p_1, p_2, \dots, p_s \rangle$ .

## 2.2 The Fundamental $m$ GFF Lemma

In this section we give the Fundamental  $m$  GFF Lemma, which is an extension of the Fundamental Lemma given by Paule. The  $\gcd(p, E^m p)$  for  $p \in K[n]$  plays a basic role in finding  $m$ -hypergeometric solutions of anti-difference equation (1.3).

**Lemma 2.1.** (“Fundamental  $m$  GFF Lemma”) Given a monic polynomial  $p \in K[n]$  with  $m$  GFF-form  $\langle p_1, p_2, \dots, p_s \rangle$ . Then

$$\gcd(p, E^m p) = [p_2]_m^1 \cdot [p_3]_m^2 \cdot \dots \cdot [p_s]_m^{s-1}.$$

**Proof.** Proceeding by induction on  $s$  the case  $s = 0$  is trivial. For  $s > 0$ ,

$$\begin{aligned} \gcd(p, E^m p) &= [p_s]_m^{s-1} \cdot \gcd([p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_{s-1}]_m^{s-1} \cdot E^{(-s+1)m} p_s, E^m ([p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_{s-1}]_m^{s-1} \cdot p_s)). \\ &= [p_s]_m^{s-1} \cdot \gcd([p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_{s-1}]_m^{s-1}, E^m ([p_1]_m^1 \cdot [p_2]_m^2 \cdot \dots \cdot [p_{s-1}]_m^{s-1})). \end{aligned}$$

The first equality is obvious, the second is a consequence of  $m$  GFF3 and  $m$  GFF4 because for  $i < s$  we have

$$\gcd([p_i]_m^i, E^m p_s) = \gcd(E^{(-s+1)m} p_s, E^m [p_i]_m^i) = E^m \gcd(E^{-sm} p_s, [p_i]_m^i) = 1.$$

together with  $\gcd(E^{(-s+1)^m} p_s, E^m p_s) \mid \gcd([p_s]_m^s, E^m p_s) = 1$ . The rest of the proof follows from applying the induction hypothesis.

□

In the above lemma we see that from the  $m$  GFF-form of a polynomial  $p$  we can find the  $m$  GFF-form of  $\gcd(p, E^m p)$ .

### 3. $m$ -Hypergeometric Solutions of Anti-Difference Equations

In this section we extend Paule approach to find  $m$ -hypergeometric solutions  $h_n$  of equation (1.3). Given an  $m$ -hypergeometric term  $a_n$  and suppose that there exists an  $m$ -hypergeometric term  $h_n$  satisfying equation (1.3), then by using (1.3) we find

$$\frac{h_n}{a_n} = \frac{h_n}{h_{n+m} - h_n} = \frac{1}{\frac{h_{n+m}}{h_n} - 1}.$$

Let  $y(n) = \frac{h_n}{a_n}$ . It follows that  $y(n)$  is a rational function of  $n$ . Let  $\langle a, b \rangle$  be the reduced form of  $w(n) = \frac{a_{n+m}}{a_n}$ . Substituting  $y(n)a_n$  for  $h_n$  in (1.3) to obtain

$$a(n)y(n+m) - b(n)y(n) = b(n). \tag{3.1}$$

This means, the problem of finding  $m$ -hypergeometric solution of (1.3) is equivalent to finding a rational solution  $y(n)$  of (3.1). If a solution  $y(n) \in K(n)$  of (3.1) with the reduced form  $\langle u, v \rangle$  exists, assume we know  $v$  or a multiple  $V \in K[n]$  of  $v$ . Then equation (3.1) can be written as

$$a(n) \cdot V(n) \cdot U(n+m) - b(n) \cdot V(n+m) \cdot U(n) = b(n) \cdot V(n) \cdot V(n+m), \tag{3.2}$$

where  $U(n) = \frac{u(n) \cdot V(n)}{v(n)}$  is unknown polynomial. Hence the problem reduces to finding a polynomial solution  $U \in K[n]$  of equation (3.2). To solve (3.2) we try to find a suitable denominator polynomial  $V$  and then  $U$  can be computed as a polynomial solution of (3.2). Let

$$v_i(n) = \frac{v(n+i)}{\gcd(v, E^m v)} \quad \text{for } i \in \{0, m\}.$$

Then (3.1) is equivalent to

$$(3.3) \quad a(n) \cdot v_0(n) \cdot u(n+m) - b(n) \cdot v_m(n) \cdot u(n) = b(n) \cdot v_0(n) \cdot v_m(n) \cdot \gcd(v, E^m v).$$

From this equation we immediately get that  $v_0(n) \mid b(n)$  and that  $v_m(n) \mid a(n)$ . Let  $m$  GFF  $(v) = \langle p_1, p_2, \dots, p_s \rangle$ , by using the  $m$  GFF-concept and the Fundamental  $m$  GFF Lemma we get that

$$v_0 = \frac{v}{\gcd(v, E^m v)} = p_1 \cdot E^{-m} p_2 \dots E^{(-s+1)m} p_s \mid b(n), \quad (3.4)$$

$$v_m = \frac{E^m v}{\gcd(v, E^m v)} = E^m p_1 \cdot E^m p_2 \dots E^m p_s \mid a(n), \quad (3.5)$$

This observation gives rise to a simple algorithm for computing a multiple  $V = [P_1]_m^1 \cdot [P_2]_m^2 \cdots [P_s]_m^s$  of  $v$ .

- Straightforward conclusion

$$p_i \mid \gcd(E^{-m} a, E^{(i-1)m} b) \quad \forall i \in \{1, \dots, s\}.$$

If  $P_i = \gcd(E^{-m} a, E^{(i-1)m} b)$  then obviously  $p_i \mid P_i$ . Thus, we could take

$$m \text{ GFF}(V) = \langle P_1, P_2, \dots, P_N \rangle,$$

where  $N = \text{dis}_m(a, b) = \max \{ i \in \mathbb{N} \mid \deg \gcd(a, E^{im} b) \geq 1 \}$ . If  $N$  is not defined then we set  $V = 1$ .

- Refined conclusions

$$p_1 \mid \gcd(E^{-m} a, b).$$

If  $P_1 = \gcd(E^{-m} a, b)$  then  $p_1 \mid P_1$  and

$$p_2 \mid \gcd\left(E^{-m} \left(\frac{a}{E^m(P_1)}\right), E^m \left(\frac{b}{P_1}\right)\right).$$

If  $P_2 = \gcd\left(E^{-m} \left(\frac{a}{E^m(P_1)}\right), E^m \left(\frac{b}{P_1}\right)\right)$ , then  $p_{12} \mid P_2$  and so on until we arrive at a  $P_N$  and we may again take  $m$  GFF  $(V) = \langle P_1, P_2, \dots, P_N \rangle$ .

□

The algorithm that we have just derived for (1.3) can be written, by using the “redefined conclusions”, as follows:

**Algorithm 3.1.**

INPUT :  $w(n) \in K(n)$  such that  $\frac{a_{n+m}}{a_n} = w(n)$  for all  $n \in \mathbb{N}$ .

OUTPUT: an  $m$ -hypergeometric solution  $h_n$  of (1.3) if it exists, otherwise “no  $m$ -hypergeometric solution of (1.3) exists”.

(1) Decompose  $w(n)$  into the reduced form  $\langle a, b \rangle$ .

(2) Compute  $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd(a, E^{im}b) \geq 1\}$ .

If  $N > 0$  then compute for  $j$  from 1 to  $N$

$$P_j(n) = \gcd(E^{-m}a, E^{(j-1)m}b)$$

$$a = \frac{a}{E^m P_j(n)}$$

$$b = \frac{b}{E^{-(j-1)m} P_j(n)}$$

$$m \text{ GFF}(V) = \langle P_1, P_2, \dots, P_N \rangle$$

otherwise  $V = 1$ .

(3) If equation (3.2) can be solved for  $U \in K[n]$  then return  $h_n = \frac{U(n)}{V(n)} a_n$ , otherwise return “no  $m$ -hypergeometric solution of (1.3) exists”.

□

#### 4. $q$ - $m$ -Greatest Factorial Factorization

In this section we define the “ $q$ - $m$ -Greatest Factorial Factorization” ( $qm$  GFF) of a polynomial which is an extension of the  $q$  GFF-concept introduced by Paule and Riese. Also, it is a  $q$ -analogue of the  $m$  GFF-concept, defined with respect to the  $q$ -shift operator  $\varepsilon$  instead of the shift operator  $E$  as for  $m$  GFF. In Sections 4 and 5, we will use  $n$  as an abbreviation for  $q^k$ .

Let us define the dispersion  $\text{dis}_m(a(n), b(n))$  of the  $q$ -monic polynomials  $a(n), b(n) \in F[n]$  is the greatest nonnegative integer  $i$  (if it exists) such that  $a(n)$  and  $b(q^{mi}n)$  have a nontrivial common divisor, i.e.,

$$\text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd(a(n), b(q^{mi}n)) \geq 1\}.$$

A polynomial  $p \in F[n]$  is said to be  $q$ -monic if  $p(0) = 1$ . Any polynomial  $p \in F[n]$  has a unique factorization, the  $q$ -monic decomposition, in the form

$$p = z \cdot n^\alpha \cdot \hat{p},$$

where  $z \in F$ ,  $\alpha \in \mathbb{N}$ , and  $\hat{p} \in F[n]$   $q$ -monic. We will write  $\gcd_q$  instead of “gcd”, indicating that the  $\gcd_q$  of two  $q$ -monic polynomials is understood to be  $q$ -monic.

More generally, if  $p_1 = z_1 \cdot n^{\alpha_1} \cdot \hat{p}_1$  and  $p_2 = z_2 \cdot n^{\alpha_2} \cdot \hat{p}_2$  are  $q$ -monic decompositions of  $p_1, p_2 \in F[n]$ , we define

$$\gcd_q(p_1, p_2) = \gcd(n^{\alpha_1}, n^{\alpha_2}) \cdot \gcd_q(\hat{p}_1, \hat{p}_2).$$

#### 4.1 Basic Definitions

**Definition 4.1.** For any  $q$ -monic polynomial  $p \in F[n]$  and  $i \in \mathbb{N}$ , the  $i$ -th  $m$ -falling  $q$ -factorial  $[p]_{m_q}^i$  of  $p$  is defined as

$$[p]_{m_q}^i = p \cdot \varepsilon^{-m} p \cdot \varepsilon^{-2m} p \cdots \varepsilon^{-(i-1)m} p.$$

For  $i = 0$ , we let  $[p]_{m_q}^0 = 1$ .

**Definition 4.2.** We say that  $\langle p_1, p_2, \dots, p_s \rangle$ ,  $p_i \in F[n]$ , is an  $qm$  GFF-form of a  $q$ -monic polynomial  $p \in F[n]$  if the following conditions hold:

$$(qm \text{ GFF1}) \quad p = [p_1]_{m_q}^1 \cdot [p_2]_{m_q}^2 \cdots [p_s]_{m_q}^s,$$

$$(qm \text{ GFF2}) \quad \text{each } p_i \text{ is } q\text{-monic and } s > 0 \text{ implies } \deg(p_s) > 0,$$

$$(qm \text{ GFF3}) \quad \gcd_q([p_i]_{m_q}^i, \varepsilon^m p_j) = 1 \text{ for } 1 \leq i \leq j \leq s,$$

$$(qm \text{ GFF4}) \quad \gcd_q([p_i]_{m_q}^i, \varepsilon^{-jm} p_j) = 1 \text{ for } 1 \leq i \leq j \leq s.$$

If  $\langle p_1, p_2, \dots, p_s \rangle$  is the  $qm$  GFF-form of a  $q$ -monic  $p \in F[n]$  we also denote this fact for short by  $qm \text{ GFF}(p) = \langle p_1, p_2, \dots, p_s \rangle$ .



### 4.2 The Fundamental $qm$ GFF Lemma

In this section we give the Fundamental  $qm$  GFF Lemma which is an extension of the Fundamental  $q$  GFF Lemma given by Paule and Riese. In finding  $m$ -hypergeometric solutions of anti-difference equations (i.e.,  $q = 1$ ) the  $\gcd(p, E^m p)$  for  $p \in K[n]$  plays a basic role. The same is true for  $qm$ -hypergeometric solutions of  $q$ -anti-difference equations with respect to the  $q$ -shift operator  $\varepsilon$  instead of  $E$ .

**Lemma 4.1.** (“Fundamental  $qm$  GFF Lemma”) *Given a  $q$ -monic polynomial  $p \in F[n]$  with  $qm$  GFF-form  $\langle p_1, p_2, \dots, p_s \rangle$ . Then*

$$\gcd_q(p, \varepsilon^m p) = [p_2]_{m_q}^1 \cdot [p_3]_{m_q}^2 \cdots [p_s]_{m_q}^{s-1}.$$

**Proof.** Proceeding by induction on  $s$  the case  $s = 0$  is trivial. For  $s > 0$ ,

$$\begin{aligned} \gcd_q(p, \varepsilon^m p) &= [p_s]_{m_q}^{s-1} \cdot \gcd_q([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1} \cdot \varepsilon^{(-s+1)m} p_s, \varepsilon^m ([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1} \cdot p_s)). \\ &= [p_s]_{m_q}^{s-1} \cdot \gcd_q([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1}, \varepsilon^m ([p_1]_{m_q}^1 \cdots [p_{s-1}]_{m_q}^{s-1})). \end{aligned}$$

The first equality is obvious, the second is a consequence of  $qm$  GFF3 and  $qm$  GFF4 because for  $i < s$  we have

$$\gcd_q([p_i]_{m_q}^i, \varepsilon^m p_s) = \gcd_q(\varepsilon^{(-s+1)m} p_s, \varepsilon^m [p_i]_{m_q}^i) = \varepsilon^m \gcd_q(\varepsilon^{-sm} p_s, [p_i]_{m_q}^i) = 1.$$

together with  $\gcd_q(\varepsilon^{(-s+1)m} p_s, \varepsilon^m p_s) \mid \gcd_q([p_s]_{m_q}^s, \varepsilon^m p_s) = 1$ . The rest of the proof follows from applying the induction hypothesis.

□

In the above lemma we see that from the  $qm$  GFF-form of a  $q$ -monic polynomial  $p$  one directly can extract the  $qm$  GFF-form of  $\gcd_q(p, \varepsilon^m p)$ .

### 5. $qm$ -Hypergeometric Solutions of $q$ -Anti-Difference Equations

In this section we extend Paule-Riese's approach to find  $qm$ -hypergeometric solutions  $g_k$  of the  $q$ -anti-difference equation

$$g_{k+m} - g_k = f_k, \tag{5.1}$$

where  $f_k$  is a given a  $qm$ -hypergeometric term. Given a  $qm$ -hypergeometric term  $f_k$  and suppose that there exists a  $qm$ -hypergeometric term  $f_k$  satisfying equation (5.1), then by using (5.1) we find

$$\frac{g_k}{f_k} = \frac{g_k}{g_{k+m} - g_k} = \frac{1}{\frac{g_{k+m} - 1}{g_k}}.$$

Let  $\tau = \frac{g_k}{f_k}$ . It follows that  $\tau$  is a rational function over  $F$ . Substituting  $\tau \cdot f_k$  for  $g_k$  in (5.1) to obtain

$$\rho \cdot \varepsilon^m \tau - \tau = 1, \quad (5.2)$$

where  $\rho = \frac{f_{k+m}}{f_k} \in F(n)$  is a rational function. Let  $\rho = z \cdot n^\alpha \cdot \frac{a}{b}$  with  $z \in F$ ,  $\alpha$  integer, and  $a, b \in F[n]$  relatively prime and  $q$ -monic. For any integer  $\alpha$  we define  $\alpha_+ = \max(\alpha, 0)$  and  $\alpha_- = \max(-\alpha, 0)$ , thus equation (5.2) is equivalent to

$$z \cdot n^{\alpha_+} \cdot a \cdot \varepsilon^m \tau - n^{\alpha_-} \cdot b \cdot \tau = n^{\alpha_-} \cdot b. \quad (5.3)$$

This means the problem of finding a  $qm$ -hypergeometric solutions of (5.1) is equivalent to finding rational solutions  $\tau$  of (5.3). Let  $\tau = \frac{u}{v}$  where  $u, v \in F[n]$  be two unknown relatively prime polynomials with  $v = n^\beta \cdot \hat{v}$  the  $q$ -monic decomposition of  $v$ . If a solution  $\tau$  of (5.3) exists, assume we know  $v$  or a multiple  $V \in F[n]$  of  $v$ . Then equation (5.3) can be written as

$$z \cdot n^{\alpha_+} \cdot a \cdot V \cdot \varepsilon^m U - n^{\alpha_-} \cdot b \cdot \varepsilon^m V \cdot U = n^{\alpha_-} \cdot b \cdot V \cdot \varepsilon^m V. \quad (5.4)$$

where  $U(n) = u(n) \cdot \frac{V(n)}{v(n)}$  is unknown polynomial. Hence the problem reduces to finding a polynomial solution  $U \in F[n]$  of equation (5.4). To solve (5.4) we try to find a suitable denominator polynomial  $V$  and then  $U$  can be computed as a polynomial solution of (5.4). Let

$$v_i = \frac{\varepsilon^i v}{\gcd_q(v, \varepsilon^m v)} \quad \text{for } i \in \{0, m\}.$$

Then (5.3) is equivalent to

$$z \cdot n^{\alpha_+} \cdot a \cdot v_0 \cdot \varepsilon^m u - n^{\alpha_-} \cdot b \cdot v_m \cdot u = n^{\alpha_-} \cdot b \cdot v_0 \cdot v_m \cdot \gcd_q(v, \varepsilon^m v). \quad (5.5)$$

From this equation we immediately get that  $v_0 \mid b$  and that  $v_m \mid a$ . Let  $qm$  GFF  $(\hat{v}) = \langle p_1, p_2, \dots, p_s \rangle$ , by using the  $qm$  GFF-concept and the Fundamental  $qm$  GFF Lemma we get that

$$v_0 = \frac{v}{\gcd_q(v, \varepsilon^m v)} = \frac{\hat{v}}{\gcd_q(\hat{v}, \varepsilon^m \hat{v})} = p_1 \cdot \varepsilon^{-m} p_2 \dots \varepsilon^{(-s+1)m} p_s \mid b, \quad (5.6)$$

$$v_m = \frac{\varepsilon^m v}{\gcd_q(v, \varepsilon^m v)} = \frac{q^{m\beta} \cdot \varepsilon^m \hat{v}}{\gcd_q(\hat{v}, \varepsilon^m \hat{v})} = q^{m\beta} \cdot \varepsilon^m p_1 \cdot \varepsilon^m p_2 \dots \varepsilon^m p_s \mid a, \tag{5.7}$$

This observation give rise to a simple algorithm for computing a multiple  $\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_s]_{m_q}^s$  of  $\hat{v}$ .

- Straightforward conclusion

$$p_i \mid \gcd_q(\varepsilon^{-m} a, \varepsilon^{(i-1)m} b) \quad \forall i \in \{1, \dots, s\}.$$

If  $P_i = \gcd_q(\varepsilon^{-m} a, \varepsilon^{(i-1)m} b)$  then obviously  $p_i \mid P_i$ . Thus, we could take

$$\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_N]_{m_q}^N,$$

where  $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd_q(a, \varepsilon^{im} b) \geq 1\}$ . If  $N$  is not defined then we set  $\hat{V} = 1$ .

- Refined conclusions

$$p_1 \mid \gcd_q(\varepsilon^{-m} a, b).$$

if we take  $P_1 = \gcd_q(\varepsilon^{-m} a, b)$  then  $p_1 \mid P_1$  and

$$p_2 \mid \gcd_q\left(\varepsilon^{-m} \left(\frac{a}{\varepsilon^m(P_1)}\right), \varepsilon^m \left(\frac{b}{P_1}\right)\right).$$

If we take  $P_2 = \gcd_q\left(\varepsilon^{-m} \left(\frac{a}{\varepsilon^m(P_1)}\right), \varepsilon^m \left(\frac{b}{P_1}\right)\right)$ , then  $p_2 \mid P_2$  and so on until

we arrive at a  $P_N$  and we may again take  $\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \dots [P_N]_{m_q}^N$ .

□

With  $\hat{V}$  in hand, all what is left for solving (5.4), and thus finding the a  $qm$ -hypergeometric solution of equation (5.1), is to determine an appropriate value of  $\gamma$  such that

$$v(n) = n^\beta \cdot \hat{v}(n) \mid V(n) = n^\gamma \cdot \hat{V}(n).$$

For that we will follow the approach given by Paule and Riese (1997). Consider equation (5.5): (i) Assume that  $\alpha \neq 0$  then either  $\alpha_- \neq 0$  or  $\alpha_+ \neq 0$ . In the first case we have  $\alpha_+ = 0$  and  $n^{\alpha_-} \mid u$ , hence  $\beta$  must be zero because of  $\gcd_q(u, v) = 1$ . This

means we can choose  $\gamma = 0$ . In the second case we have  $\alpha_- = 0$  and  $n^{\min(\alpha_+, \beta)} \mid u$ , because of  $\gcd_q(v, \varepsilon^m v) = n^\beta \cdot \gcd_q(\hat{v}, \varepsilon^m \hat{v})$ . Again  $\beta$  must be zero, and again we can choose  $\gamma = 0$ . (ii) Assume that  $\alpha = 0$ . In this case equation (5.5) evaluated at  $n = 0$  turns into

$$(z - q^{\beta m})u(0) = q^{\beta m} \cdot \delta_{0, \beta},$$

where  $\delta_{0, \beta}$  denotes the Kronecker symbol. This means if  $\beta > 0$  we obtain, observing that  $u(0) \neq 0$  in this case, as a condition for  $\beta$  that  $z = q^{\beta m}$ . Hence in case  $\alpha = 0$ , we choose  $\gamma = \frac{1}{m} \cdot \log_q(z)$  if  $z$  is a positive integer power of  $q$ , or  $\gamma = 0$  otherwise.

The algorithm that we have just derived for (5.1) can be written, by using the “redefined conclusions”, as follows:

**Algorithm 5.1.**

*INPUT* :  $\rho \in F(n)$  such that  $f_{k+m}/f_k = \rho(q^k)$  for all  $n \in \mathbb{N}$ .

*OUTPUT*: a  $qm$ -hypergeometric solution  $g_k$  of (5.1) if it exists, otherwise “no  $qm$ -hypergeometric solution  $g_k$  of (5.1) exists”.

(1) Decompose  $\rho$  into the form  $\rho = z \cdot n^\alpha \cdot \frac{a}{b}$  such that  $z \in F$ ,  $\alpha$  integer, and  $a, b \in F[n]$  relatively prime and  $q$ -monic.

(2) Compute  $N = \text{dis}_m(a, b) = \max \{i \in \mathbb{N} \mid \deg \gcd_q(a, \varepsilon^{im} b) \geq 1\}$ .

If  $N > 0$  then compute for  $j$  from 1 to  $N$

$$P_j(n) = \gcd_q(\varepsilon^{-m} a, \varepsilon^{(j-1)m} b)$$

$$a = \frac{a}{\varepsilon^m P_j(n)}$$

$$b = \frac{b}{\varepsilon^{-(j-1)m} P_j(n)}$$

$$\hat{V} = [P_1]_{m_q}^1 \cdot [P_2]_{m_q}^2 \cdots [P_N]_{m_q}^N$$

otherwise  $\hat{V} = 1$ .

(3) Determine the value of  $\gamma$  as follows:

$$\gamma = \begin{cases} \frac{1}{m} \cdot \log_q(z) & \text{if } \alpha = 0 \text{ and } z \text{ a positive integer power of } q, \\ 0 & \text{otherwise.} \end{cases}$$

(4) Take  $V = n^\gamma \cdot \hat{V}$ . If equation (5.4) can be solved for  $U \in F[n]$  then return

$g_k = \frac{U(n)}{V(n)} f_k$ , otherwise return “no  $qm$ -hypergeometric solution of (5.1) exists”. □

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